
TECHNICAL CONSIDERATIONS WHEN IMPLEMENTING LOCALTALK LINK ACCESS PROTOCOL

The LLAP Protocol is an important part of the Appletalk network system. It manages access to the node-to-node transmission of network data packets, governs access to the link, and provides a means for nodes to discover valid addresses...all error free.

INTRODUCTION

The LLAP (LocalTalk Link Access Protocol) is the ISO/OSI (International Standards Organization/Open Systems Interconnection) link layer protocol of the AppleTalk network system. This protocol manages the node-to-node transmission of data packets in the network. LLAP governs access to the link and provides a means for nodes to discover valid addresses. It does not guarantee packet delivery; it does guarantee that those packets that are delivered are error-free.

This Appnote (Application Note) does not address the architectural issues of writing a driver but it does focus on the details of using an SCC to send and receive LLAP frames. However, some of the problems of transmitting and receiving LLAP frames are discussed, using sample code written for Zilog's Z80181 Emulation Adapter Board. Also, the problems of sending sync pulses, timing transmissions and determining that a frame has been received properly will be discussed.

GENERAL DESCRIPTION

The LocalTalk Link Access Protocol (LLAP) is the ISO-OSI link layer protocol of the AppleTalk network system using LocalTalk. Along with ELAP (the corresponding Ethernet link layer protocol) and TLAP (the Token Ring link layer protocol), it provides the foundations upon which the other protocols rest. The LLAP protocol supports the node-to-node transmission of packets used by DDP and RTMP to route packets around the internetwork; DDP, in turn, supports the name binding functions of NBP, the reliable frame delivery of ATP, and the rest of the AppleTalk protocol stack.

A majority of the difficult timing and all of the hardware interface problems crop up in the LLAP driver. These problems are so difficult that it makes sense to start writing such a driver by writing experimental routines that transmit and receive frames. This App Note addresses the intricacies of the interframe and interdialog timings before trying to engineer code that will truly be a driver. Also, some of the experimental routines to run on the Z80181 Emulation Adapter Board will be explained.

The LLAP provides the basic transmission of packets from one node to another on the same network. LLAP accepts packets of data from clients residing on a particular node and encapsulates that data into its proper LLAP data packet. The encapsulation includes source and destination addresses for proper delivery. LLAP ensures that any damaged packet is discarded and rejected by the destination node. The LLAP makes no effort to deliver damaged packets.

Carrier Sense Multiple Access with Collision Avoidance

It is LLAP's responsibility to provide proper link access management to ensure fair access to the link by all nodes on that network. The access discipline that governs this is known as Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA). A node wishing to gain access to the link must first sense that the link is not in use by any other node (carrier sense); if the link has activity, then the node wishing to transmit must defer transmission. The ability of LLAP to allow multiple access to the link also

GENERAL DESCRIPTION (Continued)

leaves room for possible collisions with other data packets. LLAP attempts to minimize this probability (collision avoidance).

Two techniques are used by LLAP in its implementations of CSMA/CA. LLAP outlines this procedure but falls short in endorsing which hardware to use. (The SCC is, of course, used by Apple.) The first technique takes advantage of the distinctive 01111110 flag bytes that encapsulate the data packet (note that this implies that SDLC is used). LLAP stipulates that a minimum of two flags precede each of these data packets. The leading flag characters provide byte synchronization and give a clue to any listener that some other node is using the link at a particular time (use the Hunt bit in RR0 if the SCC is used).

In SDLC mode, the receiver automatically synchronizes on the flag byte and resets the Hunt bit to zero. The SCC has some latency in detecting these flag bytes due to the shifter, etc. This is not ideal because the node needing to transmit may determine that the link is free, when in fact the flag bytes are still being shifted into its receiver (i.e., the link is not idle at all).

A closing flag is also needed to fully encapsulate the data packet. LLAP requires that 12 to 18 ones be sent after this closing flag. The LocalTalk hardware (i.e., the SCC) interprets this as an abort sequence and causes the node's hardware to lose byte sync; this then confirms that the current sender's transmission is over. In SDLC mode, seven or more contiguous 1's in the receive data stream forces the receiver into Hunt (Hunt bit set) and an

External/Status interrupt can be generated. This is important because the node wishing to use the bus can simply wait for this interrupt before preparing to transmit its packet.

LLAP uses a second technique in its carrier sensing. LLAP requires that a synchronization pulse for an idle period of at least two bit times be transmitted prior to sending the RTS handshaking frame (Figure 1). This synchronization is obtained by first enabling the hardware line so that an edge is detected by all the receivers on the network. This initial edge is perceived as the beginning of the clocking period. It is soon followed by an idle period (a period with no carrier) of at least two bit times. All the receivers on the network see this idle period and assume that the clock has been lost (missing clock bit set on RR10). This method is much more immediate than the byte flag synchronization method and provides a quicker way of determining whether the link is in use. Unfortunately, an interrupt is not generated by this missing clock and, therefore, polling must be implemented.

The Z80181 code used for polling the missing clock bit is approximately fifty clock cycles which at 10 MHz is about 5 μ sec or about one bit time. This is still relatively quicker than the time required for the SCC to reset the Hunt bit (the flag character takes at least eight bit times for it to be shifted through the buffer before the Hunt bit is reset to zero). Synchronization pulses can be sent before every frame but because of the time constraints associated with the interframe gaps it makes sense to send such pulses only before the lapENQ and lapRTS frames.

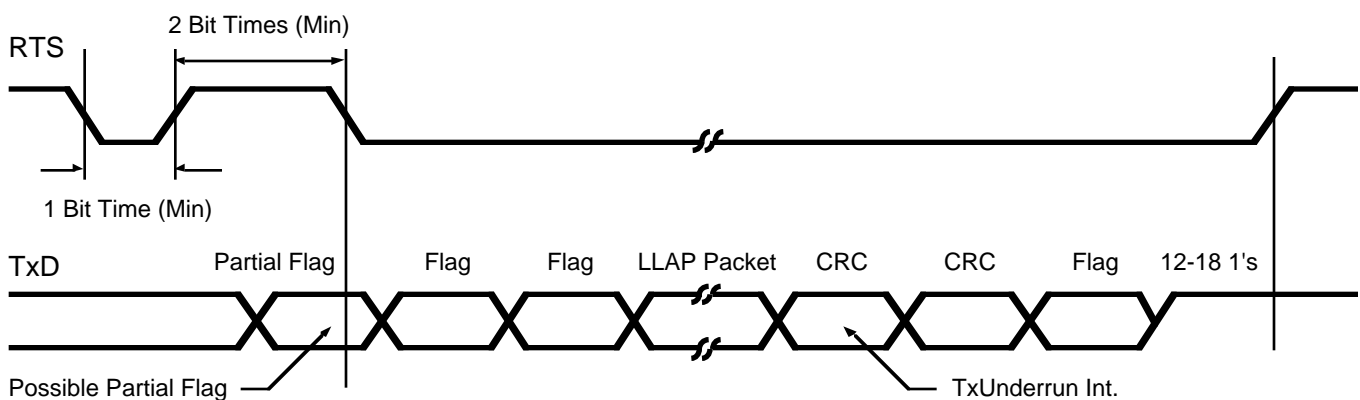


Figure 1. CSMA/CA Synchronization Pulse Timing Diagram

Dynamic Node ID

LLAP requires the use of an 8-bit node identifier number (node ID) for each node on the link. Apple had decided that all LLAP nodes must have a dynamically assigned node ID. A node would assign itself its unique address upon activation. It is then up to that particular node to ascertain that the address it had chosen is unique. A node randomly chooses an 8-bit address (for example, the refresh register value on the Z80181 is added to a randomly chosen value on the receive buffer to obtain a pseudo random 8-bit address).

The node then sends out an LLAP Enquiry control packet to all the other nodes and waits for the prescribed interframe gap of 200 μ sec. If another node is already using this node ID, then that node must respond within 200 μ sec with a LLAP Acknowledgment control packet. The new node must then rebroadcast a new guess for its node ID. If a LLAP Acknowledgment packet is not received within 200 μ sec then the new node assumes that the address is indeed unique. The new node must rebroadcast the LLAP enquiry packet several more times to account for cases when the packet could have been lost or when the guessed node ID is busy and could have missed the Enquiry packet.

LLAP Packet

LLAP packets are made up of three header bytes (destination ID, source ID and LLAP type) and 0 to 600 bytes of variable length data. The LLAP type indicates the type of packet that is being sent. 80H to FFH are reserved as LLAP control packets. The four LLAP control packets that are currently being used are: The lapENQ, which is used as enquiry packet for dynamic node assignments; the lapACK, which is the acknowledgment to the lapENQ; the lapRTS, which is the request to send packet that notifies the destination of a pending transmission; and the lapCTS, which is the clear-to-send packet in response to the RTS packet. Control packets do not contain data fields.

LLAP Packet Transmission

LLAP distinguishes between two types of transmissions: a directed packet is sent from the source node to a specific destination node through a directed transmission dialog; a broadcast packet is sent from the source node to all nodes on the link (destination ID is FFH) through a broadcast transmission dialog. All dialogs must be separated by a minimum Inter Dialog Gap (IDG) of 400 μ sec. Frames within these dialogs must be separated from each other with a maximum Inter Frame Gap (IFG) of 200 μ sec.

The source node uses the physical layer to detect the presence or the absence of data packets on the link. The node will wait until the line is no longer busy before attempting to send its packets. If the node senses that the line is indeed busy, then this node must defer. When the node senses that the line is idle, then the node waits the minimum IDG plus some randomly generated time before sending the packet (the line must remain idle throughout this period before attempting to send the packet). The initial packets to be sent are handshaking packets. The first packet sent by the source node to its destination node is the RTS packet. The receiver of this RTS packet must return a CTS packet within the allowable maximum IFG. The source node then starts transmitting the rest of its data packet upon receiving this CTS.

Collisions are more likely to occur during the handshaking phase of the dialog. The randomly generated time that is added to the IDG tends to help spread out the use of the link among all the transmitters. A successful RTS to CTS handshake signifies that a collision did not take place. An RTS packet that collides with another frame has corrupt data that shows up as a CRC error on the receiving or the destination node. Upon receiving this, the destination node infers that a collision must have taken place and abstains from sending its CTS packet. The source or the transmitting node sees that the CTS packet was not received during the IFG and also infers that a collision did take place. The sending node then backs off and retries.

The LLAP keeps two history bytes that log the number of deferrals and collisions during a dialog. These history bytes help determine the randomly generated time that is added to the IDG. The randomly generated time is readjusted according to the traffic conditions that are present on the link. If collisions or deferrals have just occurred on the most recently sent packets, then it can be assumed that the link has heavier than usual traffic. Here, the randomly generated number should be a larger number in order to help spread out the transmission attempts. Similarly, if the traffic is not so great, then the randomly generated number should be smaller, thus reducing the dispersion of the transmission attempts.

LocalTalk Physical Layer

LocalTalk uses the SDLC format and the FM0 bit encoding technique. The RS-422 signalling standard for transmission and reception was chosen over the RS-232 because a higher data rate over a longer physical distance is required. LocalTalk requires signals at 230.4 Kbits per second over a distance of 300 meters.

HARDWARE CONFIGURATION

As shown in Figure 2, the hardware used to implement this LLAP driver consists of the Z80181 (an integration of the Z80180 compatible MPU core with one channel of a Z85C30 SCC, Z80 CTC, two 8-bit general-purpose parallel ports and two chip select signals) operating at 10 MHz, a 3.6864 MHz clock source and an RS-422 line driver with tri-state.

The SCC's clocking scheme decouples the micro-processor's clock from the communication clock (Figure 3). The DPLL uses the /RTxC pin as its source. The /RTxC also drives the Baud Rate Generator which divides its input by sixteen. The resulting 230.4 kHz signal is then

used as transmitter clock. This 230.4 kHz signal is also used by one of the Z80181's counter/timer trigger inputs (Z80 CTC's channel 1) which is used to count the number of elapsed bit times. In counter mode, each active edge to the CTC's CLK/TRG1 input causes the downcounter of the CTC to be decremented. The /TRxC pin is programmed as BRG output and is connected to the CLK/TRG1 input through an external wire.

The /RTS signal is used to tri-state RS-422 to allow other node transmitters to drive the line. This signal is asserted and deasserted through bit1 of the SCC's Write Register 5.

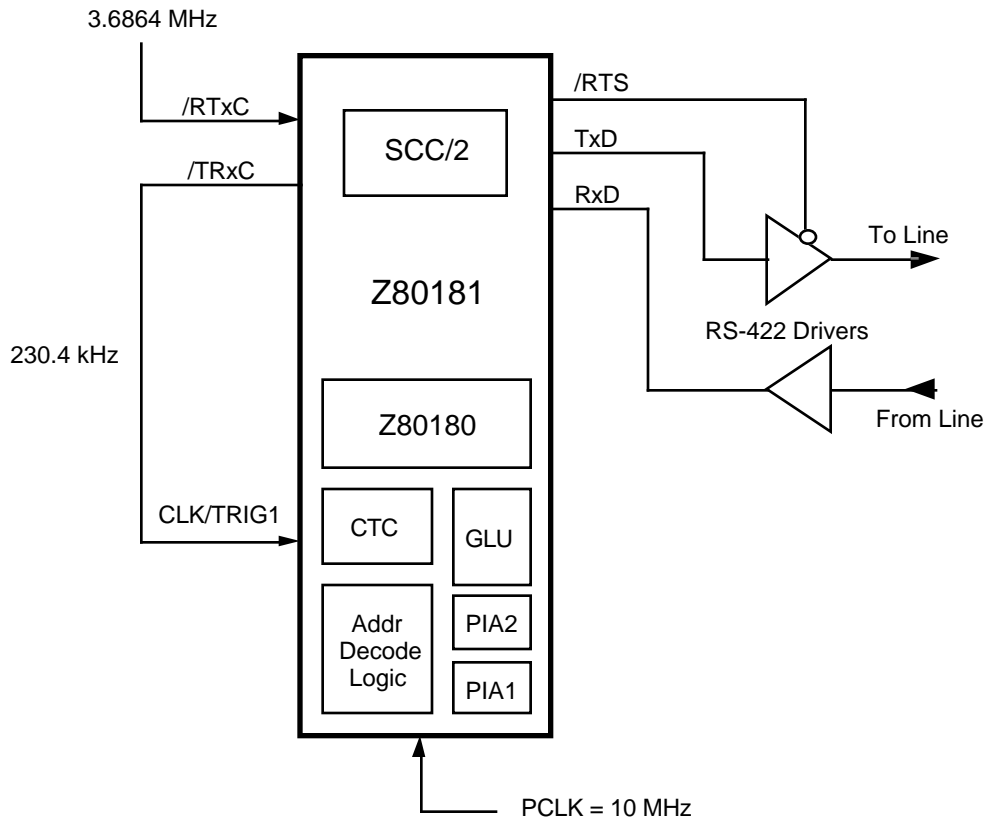


Figure 2. Driver Hardware Configuration

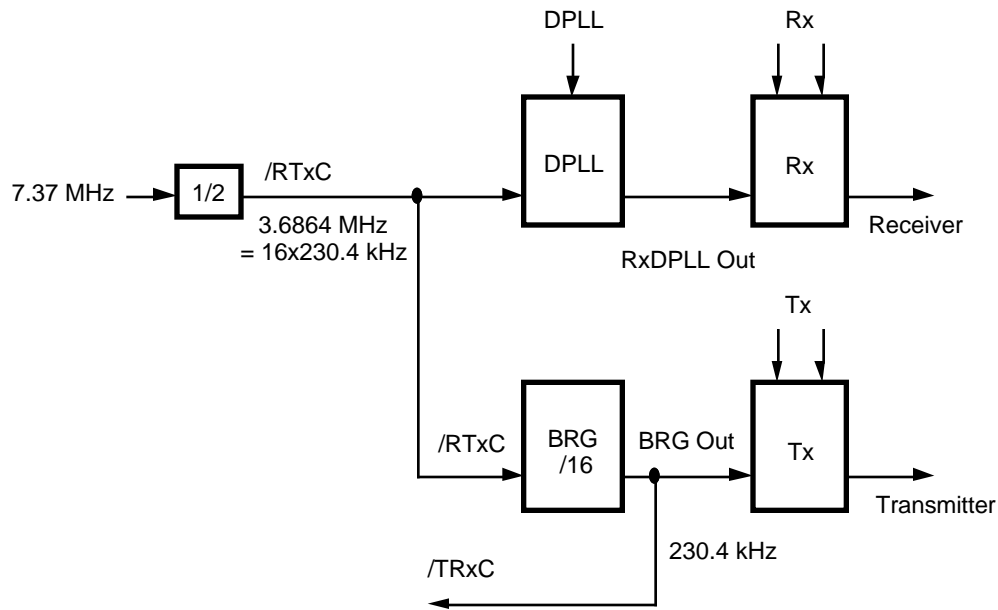


Figure 3. SCC Clocking Scheme

Listing 1 (Reference Appendix A for Listings 1 through 4) shows the assembler code for this SCC initialization. Note that the SCC is treated as a peripheral by the Z80181's MPU. For example, an I/O write to the `scc_cont` (address e8H) or to the `scc_data` (address e9H) is a write to the SCC's control and data registers, respectively. As shown in Listing 1, the SCC is initialized by issuing I/O writes to the pointer and then to the control registers in an alternating fashion. It is therefore very important that all interrupts are disabled during this initialization routine.

The SCC is initially reset through software before proceeding to program the other write registers. A NOP is sufficient to provide the four PCLKs required by the SCC recovery time after a soft reset. The SCC is programmed for SDLC mode. The receive, transmit and external interrupts are all initially disabled during this initialization. Each of these interrupt sources are enabled at their proper times in the main program. The SCC is programmed to include status information in the vector that it places on the bus in response to an interrupt acknowledge cycle (see Listing 4 of the SCC interrupt vector table for all the possible sources).

Since SDLC is bit-oriented, the transmitter and receiver are both programmed for 8 bits per character as required by LLAP. Address filtering is implemented by setting the Address Search Mode bit 2 on WR3. Setting this bit causes messages with addresses not matching the address programmed in WR6 and not matching the broadcast address to be rejected. Values in WR10 presets the CRCs to ones, sets the encoding to FM0 mode and makes certain that transmission of flags occur during idle and underrun conditions. WR11 is set so that the receive clock is sourced by the DPLL output; the transmit clock is sourced by the Baud Rate Generator output; /TRxC's output is from the BRG. The input to the BRG is from the /RTxC.

The BRG's time constant is loaded in WR13 and WR12 so that the /RTxC's 3.6864 MHz signal is divided by 16 in order to obtain a 230.4 kHz signal for the transmitter clock. WR14 makes certain that the DPLL is disabled before choosing the clock source and operating mode. The DPLL is enabled by issuing the Enter Search Mode in WR14.

TRANSMITTING A LLAP FRAME

Listing 2 shows the assembler code for subroutine txenq, which sends an lapENQ frame on the line once the system has determined that the line is quiet. Note that this routine can easily be generalized to send any frame.

The first responsibility of txenq is to send the sync pulse required by the CSMA/CA protocol. To do this, txenq sets the /RTS pin active low, enabling the transmitter drivers, and then sets it high again to disable them. In order to satisfy the requirements of the CSMA/CA protocol, the transmitter drivers must remain off for at least one bit time (4.3 μ sec) to guarantee that all the receivers see at least one transition. Our routine satisfies this requirement because the two ld instructions (7 T states each), the two nop instructions (4 T states each) and the two "out" instructions (11 T states each) required to set the /RTS line high, take more than 4.3 μ sec to execute on the 10 MHz Z80181. The transmitter drivers must then remain off for at least two bit times in order to ensure that all receivers lose clock; again, the routine depends upon the time required to execute the instructions before we turn the transmitter drivers on again.

After sending the sync pulse and waiting the required period of silence, txenq turns on the transmitter drivers to send the frame. Now, the routine must wait while the SCC sends out the leading flags. This takes 16 bit times, and since the SCC does not tell when this has happened, the transmit routine has no choice but to time this. Our routine does this by calling bit time, which is discussed below.

When the two flags have been transmitted, the first data byte is written to the data register of the SCC. Thereafter, the routine polls the SCC status register, and when that register shows the transmit buffer register is empty, the routine sends the next data character. This polling method can obviously be replaced by an interrupt routine that is entered when the transmit buffer is empty or by setting up the Z80181's DMA to send characters on demand to the SCC.

After the first data byte is transmitted, the txenq routine sets the SCC to mark on idle so that the abort is sent when the frame is over. This operation can be done any time after the first data character has been placed in the transmit buffer and before the trailing flag is shifted out. Txenq asserts this mark on idle command after the first character is placed in the transmit buffer so that LLAP has control and that no issues of latency may arise (particularly in designs using interrupt or DMA).

After the last data byte is written to the SCC, the transmit routine must wait while the last data byte (the one that the SCC had just sent to shifter), the two CRC bytes, one flag byte and 12 to 18 bit times of marking are transmitted. This total of 44 to 50 bit times is again timed by bittime. When bittime indicates that enough time has elapsed, the transmitter drivers are turned off.

Since our hardware includes connecting the output of the baud rate generator to the input of counter/timer 1 on the Z80181, that counter timer counts the bit times. The bit time routine feeds an appropriate count value into the counter and enables an interrupt routine to receive control when the count expires. The interrupt routine ctc1int, shown in Listing 4, sets the timeflag which the transmit routine polls.

Note that the call to bittime, the interrupt routine, the polling code and the length of time it takes to write to the SCC registers after a polling loop is exited, all take up a time that can be a significant number of bits. In order to do these timings accurately, calculate the number of PCLK cycles required to execute these pieces of code and to adjust the counter value that bittime requires. This adjustment is dependent on the hardware configuration and on the exact implementation details of the code.

Note, incidentally, that software must put the entire frame into the transmit register, including the addresses. The SCC does not generate addresses even when set in WR6.

RECEIVING LLAP FRAMES

In the experiments, the interrupt routines were used to receive characters and to handle special conditions when the frame is complete. Listing 3 shows the interrupt handlers that are entered when the SCC receives a character and when the SCC interrupts for a special condition (typically, end of frame). As with transmission, it is obvious that we could receive characters by polling the SCC (using up all available CPU cycles) or by using DMA (using up very few). It is estimated that the recint routine uses up about 1/3 of the available 34 μ sec (4.3 μ sec x 8-bit times) cycles on a 10 MHz processor.

The recint routine moves each character as it is received from the SCC to a memory buffer and increments the buffer pointer. The frame's data length is checked to make certain that the maximum allowable frame size is not exceeded.

The spcond interrupt handler checks the status of the SCC to find out what has happened. The presence on an overrun condition or a CRC error is flagged as a receive frame error.

The second routine decrements the receiver buffer address by two to account for the two CRC bytes that are read from the SCC before the special condition interrupt occurs. Note that the SCC does not filter these CRC bytes, nor does it filter the address byte. Everything received after the leading flags and before the trailing flags appears in the receive buffer.

One difficulty that arises in LLAP that was not addressed here is that the receipt of a frame very often creates an obligation to send a frame back to the sender within the

interframe gap, which is 200 μ secs. This difficulty is even greater than it might appear. The 200 μ sec gap starts when the frame is received; it ends when the leading flags and destination address of the response are sent. Start sending the response soon enough to allow sending two leading flags (plus a possible leading flag fragment) and the first data character, and to allow for the 3-bit delay in the SCC shifter. Therefore, start sending early enough to transmit 35 bits before the interframe gap expires, or about 70 μ secs after you receive the frame.

CONCLUSIONS

The problems of sending the sync pulses, the timing of transmission packets, and the problems associated with the reception of packets as defined by LLAP are handled by the Z80181 and its peripherals. It was demonstrated that LLAP frames can be transmitted and received by using the straight forward polling method and by using interrupt routines. In a much busier environment where the processor cannot strictly be an LLAP engine, other

methods such as using DMA in a fully interrupt driven environment must be used. It was also demonstrated that severe CPU overhead is used in setting up the sync pulses, timing out delays, etc., before each LLAP frame. A modified SCC that transmits and receives special LLAP frames helps in off loading some of this overhead, hence freeing the CPU to do other tasks.

APPENDIX A

Listing 1- Asembler Code for SCC Initialization

LISTING 1

```

475
476
477
000001e2 478 initscc:
000001e2 f3 479          di
000001e3 f5 480          push     af
000001e4 c5 481          push     bc
000001e5 e5 482          push     hl
483
000001e6 3e09 484          ld      a,09h
000001e8 d3e8 485          out     (scc_cont),a
000001ea 3e80 486          ld      a,80h
000001ec d3e8 487          out     (scc_cont),a
000001ee 00 488          nop
489
490
000001ef 21Wwww 491          ld      hl,scctable
000001f2 492 scc1:
000001f2 7e 493          ld      a,(hl)
000001f3 feff 494          cp     0ffh
000001f5 caWwww 495          jp     z,finiscc
000001f8 d3e8 496          out     (scc_cont),a
000001fa 23 497          inc    hl
000001fb 7e 498          ld      a,(hl)
000001fc d3e8 499          out     (scc_cont),a
000001fe 23 500          inc    hl
000001ff c3R000+01f2, 501          jp     scc1
502
00000202 503 scctable:
00000202 04 504          db     04h
00000203 20 505          db     00100000b
506
00000204 01 507          db     01h
00000205 00 508          db     00h
509
00000206 02 510          db     02h
00000207 00 511          db     00h
512
00000208 03 513          db     03h
00000209 cc 514          db     0cch
515
516
517
0000020a 05 518          db     05h
0000020b 60 519          db     60h
520
0000020c 06 521          db     06h
0000020d 00 522          db     00h
523
0000020e 07 524          db     07h
0000020f 7e 525          db     7eh
526
00000210 09 527          db     09h
00000211 01 528          db     01h
529
530
531

```


00000212 0a	532	db	0ah	;WR10
00000213 e0	533	db	0e0h	;crc preset to one,fm0, flag idle/undr
	534			
00000214 0b	535	db	0bh	;WR11
00000215 f6	536	db	0f6h	;rtxc=xtal,rxcd=dp11,txc=brg,trxc=brg out
	537			
00000216 0c	538	db	0ch	;WR12
00000217 06	539	db	06h	;brg tc low,for 230.4kbps using rtxc=3.68MHz
	540			
00000218 0d	541	db	0dh	;WR13
00000219 00	542	db	00h	;brg tc high
	543			
0000021a 0e	544	db	0eh	;WR14
0000021b 60	545	db	60h	;disable dp11
	546			;no local loop back,brg source=rtxc
	547			
0000021c 0e	548	db	0eh	;WR14
0000021d c0	549	db	0c0h	;select fm mode
	550			;no local loop back,brg source=rtxc
	551			
0000021e 0e	552	db	0eh	;WR14
0000021f a0	553	db	0a0h	;dp11 source=rtxc,
	554			;no local loop back,brg source=rtxc
	555			
00000220 0e	556	db	0eh	;WR14
00000221 20	557	db	20h	;enter search mode
	558			;no local loopback
	559			
00000222 0e	560	db	0eh	;WR14
00000223 01	561	db	01h	;null,no local loopback,enable the brg
	562			
00000224 03	563	db	03h	;WR3
00000225 cc	564	db	0cch	;rx 8b/c,enable rx crc,addrs src,rx disable
	565			
00000226 0f	566	db	0fh	;WR15
00000227 00	567	db	00h	;ext/stat not used
	568			
	569			;WR0
00000228 10	570	db	10h	;reset ext/stat once
00000229 10	571	db	10h	;reset ext/stat twice
	572			
0000022a 01	573	db	01h	;WR1
0000022b 00	574	db	00h	;disable all int sources
	575			
0000022c 09	576	db	09h	;WR9
0000022d 09	577	db	09h	;enable int
0000022e ff	578	db	0ffh	;finished
	579			
0000022f	580	finscc:		
0000022f e1	581	pop	hl	;
00000230 c1	582	pop	bc	;
00000231 f1	583	pop	af	;
00000232 c9	584	ret		
	585			

LISTING 2

```

600                                     ;*****
601                                     ;
602                                     ;Subroutine to transmit the llapenq packet
603 txenq:                               ;*****
604                                     ;
00000244 603 txenq:
605         push        af                ;save status and a reg
00000244 f5 606         push        bc                ;save
00000245 c5 607         push        hl                ;save
00000246 e5 608                                     ;
609         di                          ;make sure that
610                                     ;no interrupt routine
611                                     ;nor should interrupt
612                                     ;occur during
613                                     ;this subroutine.
00000248 3e03 614         ld          a,03h
0000024a d3e8 615         out         (scc_cont),a    ;WR3
0000024c 3ecc 616         ld          a,0cch
0000024e d3e8 617         out         (scc_cont),a    ;8b/char,rx crc
618                                     ;enable,addrs src
619                                     ;and rx disabled
620
00000250 3e0a 621         ld          a,0ah                ;select WR10
00000252 d3e8 622         out         (scc_cont),a
00000254 3ee0 623         ld          a,11100000b        ;idle with flags
00000256 d3e8 624         out         (scc_cont),a
625
626                                     ;****enable transmitter ****
00000258 3e05 627         ld          a,05h                ;select WR5
0000025a d3e8 628         out         (scc_cont),a
0000025c 3e68 629         ld          a,01101000b        ;enable tx
0000025e d3e8 630         out         (scc_cont),a
631                                     ;
632                                     ;
633                                     ;****enable rs-422 driver ****
00000260 3e05 634         ld          a,05h                ;select WR5
00000262 d3e8 635         out         (scc_cont),a
00000264 3e6a 636         ld          a,01101010b        ;enable tx,
00000266 d3e8 637         out         (scc_cont),a        ;reset rts
00000268 00 638         nop
00000269 00 639         nop
640                                     ;nop's needed to complete 4.3 usec
641                                     ;for 1 bit time enable of transmitter.
642                                     ;total delay=2*(7+11+4) T states at 10 MHZ
643                                     ;
644                                     ;****disable rs-422 driver for 2 bit times****
0000026a 3e05 645         ld          a,05h                ;select WR5
0000026c d3e8 646         out         (scc_cont),a
0000026e 3e68 647         ld          a,01101000b        ;enable tx, set rts
00000270 d3e8 648         out         (scc_cont),a
649                                     ;
00000272 3e80 650         ld          a,10000000b        ;reset txcrc
00000274 d3e8 651         out         (scc_cont),a
00000276 0601 652         ld          b,01h                ;delay count
00000278 653 csloop:
00000278 10fe 654         djnz        csloop        ;loop needed
655                                     ;to complete
656                                     ;8.6 usec min.
657                                     ;or 2 bit times.
658                                     ;****enable rs-422 driver for llap transmission****
0000027a 3e05 659         ld          a,05h                ;select WR5
    
```

0000027c d3e8	660	out	(scc_cont),a	
0000027e 3e6b	661	ld	a,01101011b	;sdlc crc,
	662			;txcrc enable,
	663			;reset rts
00000280 d3e8	664	out	(scc_cont),a	
	665			
	666			;
	667			;**start counting out 2 flag character times **
	668			;
	669			;count 16 bit times
	670			;from the rs-422 enable
	671			;for 2 flags.
	672			;btdelay=subr delay+ctc1int+polling=8bits
	673			;16 bit times-btdelay=16-8=08h
	674			;
00000282 0e08	675	ld	c,08h	
00000284 cdWwww	676	call	bittime	;bittime delay
	677			;is stored in reg.c
	678			;and bit1 of timflg
	679			;will indicate
	680			;count termination.
	681			
00000287682 l6:				;timer flag
00000287 3aWwww	683	ld	a,(timflg)	;
0000028a cb4f	684	bit	1,a	;if bit1=1 then
	685			;count terminated
0000028c 28f9	686	jr	z,l6	;
0000028e cb8f	687	res	1,a	;reset timflg bit1
00000290 32Wwww	688	ld	(timflg),a	;update timflg
	689			;
00000293 3e03	690	ld	a,03h	;
00000295 d3e5	691	out	(ctc1_cont),a	;disable int,
	692			;software reset
	693			;to kill the counter1
00000297 0602	694	ld	b,02h	;2+1 bytes to transmitted
00000299 21Wwww	695	ld	hl,txlapenq	;point to txlapenq buffer
	696			;send 1st byte
0000029c 7e	697	ld	a,(hl)	;
0000029d d3e9	698	out	(scc_data),a	;and send it
0000029f 23	699	inc	hl	;point to the next byte
	700			;
000002a0 3ec0	701	ld	a,0c0h	;reset eom latch command
000002a2 d3e8	702	out	(scc_cont),a	
	703			;
000002a4 f3	704	di		;disable all int
000002a5 3e0a	705	ld	a,0ah	;select WR10
000002a7 d3e8	706	out	(scc_cont),a	
	707			;idle with 1's
	708			;at the end of the frame
000002a9 3ee8	709	ld	a,11101000b	
000002ab d3e8	710	out	(scc_cont),a	
	711			;
000002ad 3e00	712 txq2:	ld	a,00h	
000002af d3e8	713	out	(scc_cont),a	;rr0
000002b1 dbe8	714	in	a,(scc_cont)	;read rr0
000002b3 cb57	715	bit	2,a	;read tx buffer empty
000002b5 28f6	716	jr	z,txq2	;loop if zero
000002b7717 txq1:				
000002b7 7e	718	ld	a,(hl)	;
000002b8 d3e9	719	out	(scc_data),a	;and send it
000002ba 23	720	inc	hl	;point to the next byte
	721			


```

784                                     ;subroutine to time out bit time 4.3 usec per bit
785                                     ;register c contains the number of bits to be
                                     ;counted down
786                                     ;*****
000002ef 787 bittime:
000002ef f5 788         push     af         ;save status and a reg
000002f0 c5 789         push     bc         ;save
000002f1 e5 790         push     hl         ;save
791                                     ;
000002f2 3ed2 792         ld      a,0d2h      ;ctc1 int vector
000002f4 d3e5 793         out     (ctc1_cont),a
794                                     ;
000002f6 3edf 795         ld      a,11011111b   ;
000002f8 d3e5 796         out     (ctc1_cont),a ;enable int
797                                     ;select counter mode
798                                     ;clk/trg edge starts with rising edg
799                                     ;time constant follows
800                                     ;software reset
000002fa 79 801         ld      a,c         ;reg c contains the number of bits
000002fb d3e5 802         out     (ctc1_cont),a ;load the number of bits to be counted
803                                     ;**
000002fd e1 804         pop     hl         ;restore
000002fe c1 805         pop     bc         ;restore
000002ff f1 806         pop     af         ;restate status and a reg
00000300 fb 807         ei
00000301 c9 808         ret
809

```

LISTING 3

```

1131      .*****
1132      ;
1133      ;receive int service routine.
1134      .*****
1135      ;save received character in receiver buffer
1136      ;to by rxpointer
0000044d 1137recint:
0000044d f5 1138      push      af          ;save af
0000044e d5 1139      push      de
0000044f e5 1140      push      hl
00000450 dbe9 1141      in        a,(scc_data) ;read scc data
00000452 2aWwww 1142      ld        hl,(rxpointer) ;
00000455 77 1143      ld        (hl),a        ;save it
00000456 23 1144      inc      hl            ;update pointer
00000457 22Wwww 1145      ld        (rxpointer),hl ;
0000045a ed5bWwww 1146      ld        de,(rxbufend) ;end of rx buffer
0000045e af 1147      xor      a            ;reset cy
0000045f ed52 1148      sbc     hl,de         ;
00000461 c2Wwww 1149      jp      nz,recexit    ;if not zero,then receive
                                byte length is ok
00000464 21Wwww 1150      ld        hl,recerrflg ;
00000467 cbc6 1151      set     0,(hl)        ;set bit0=1 maxfrmflg to indicate error
                                ;because of max frame
                                size exceeded.
00000469 1153recexit:
00000469 3e38 1154      ld        a,038h
0000046b d3e8 1155      out     (scc_cont),a  ;reset highest ius
0000046d e1 1156      pop     hl
0000046e d1 1157      pop     de
0000046f f1 1158      pop     af            ;restore af
00000470 fb 1159      ei      ;enable int
00000471 c9 1160      ret     ;return from int
1161      ;note ret and not reti is used for scc
1162      ;interrupts on the z80181.
1163
1164      .*****
1165      ;special receive interrupt service routine
1166      .*****
1167      ;
1168      ;"parity is special condition" bit is off.
1169      ;special conditions are eof or rx overrun error.
1170      ;crc error flag is valid only if eof is valid.
1171      ;if frame is ok then recerrflg bit1=0, otherwise
00000472 1172 spcond:
00000472 f5 1173      push     af          ;save af reg
00000473 c5 1174      push     bc
00000474 e5 1175      push     hl;
1176
00000475 3e01 1177      ld        a,01h
00000477 d3e8 1178      out     (scc_cont),a  ;read rr1
00000479 dbe8 1179      in        a,(scc_cont)
0000047b e660 1180      and     01100000b    ;check bit6 (crc) or bit5 (overrun)
0000047d caWwww 1181      jp      z,ok
1182      ;
00000480 21Wwww 1183      ld        hl,recerrflg ;fetch receive error flag
00000483 cbce 1184      set     1,(hl)        ;set bit1=1 for frame not ok
00000485 c3Wwww 1185      jp      crc_exit
    
```

```

00000488          1186      ok:
00000488 21Wwww   1187      ld      hl,recerrflg      ;fetch receive error flag
0000048b cb8      1188      res      1,(hl)          ;set bit1=0 for frame ok
                  1189
                  1190
0000048d          1191  crc_exit:
0000048d dbe9      1192      in      a,(scc_data)    ;read 2nd crc (debug only) and
0000048f 2aWwww   1193      ld      hl,(rxpointer) ;load pointer
00000492 2b       1194      dec     hl              ;adjust rx buff ptr for crc1
00000493 2b       1195      dec     hl              ;adjust rx buff ptr for crc2
00000494 22Wwww   1196      ld      (rxpointer),hl ;
                  1197
00000497          1198  spexit:
00000497 3e38      1199      ld      a,038h
00000499 d3e8      1200      out     (scc_cont),a   ;reset highest ius
                  1201
0000049b e1       1202      pop     hl              ;restore hl
0000049c c1       1203      pop     bc              ;restore bc
0000049d f1       1204      pop     af              ;restore af
0000049e fb       1205      ei      ;enable int
0000049f c9       1206      ret     ;return from int
                  1207

```

LISTING 4

```

1306
1307
1308
1309
00000509 1310 ctc1int:
1311
1312
1313
1314
00000509 f5 1315 push af
0000050a c5 1316 push bc
0000050b e5 1317 push hl
1318
0000050c 21Wwww 1319 ld hl,timflg
0000050f 7e 1320 ld a,(hl)
00000510 cbcf 1321 set 1,a
00000512 77 1322 ld (hl),a
00000513 e1 1323 pop hl
00000514 c1 1324 pop bc
00000515 f1 1325 pop af
00000516 fb 1326 ei
00000517 ed4d 1327 reti
1328
1329
1330
1331
1332
1333
1334
1335
1336
00000a00 1337 org sdlc + 0a00h
00000a00 1338 sccvect:
1339 if scc_a
00000a00 1340 .block 8
1341 .endif
00000a08 R000+03e9, 1342 dw txint
00000a0a R000+04c8, 1343 dw ext_stat
00000a0c R000+0433, 1344 dw recint
00000a0e R000+0454, 1345 dw spcond
1346
1347 if not scc_a
00000a10 1348 .block 8
1349 .endif
1350
00000a18 1351 temp: .block 1
1352
1353
1354
1355
1356
0000ad0 1357 org 0ad0h
0000ad0 R000+04d8, 1358 dw ctc0int
0000ad2 1359 org 0ad2h
0000ad2 R000+0509, 1360 dw ctc1int
1361
1362
1363
1364

```

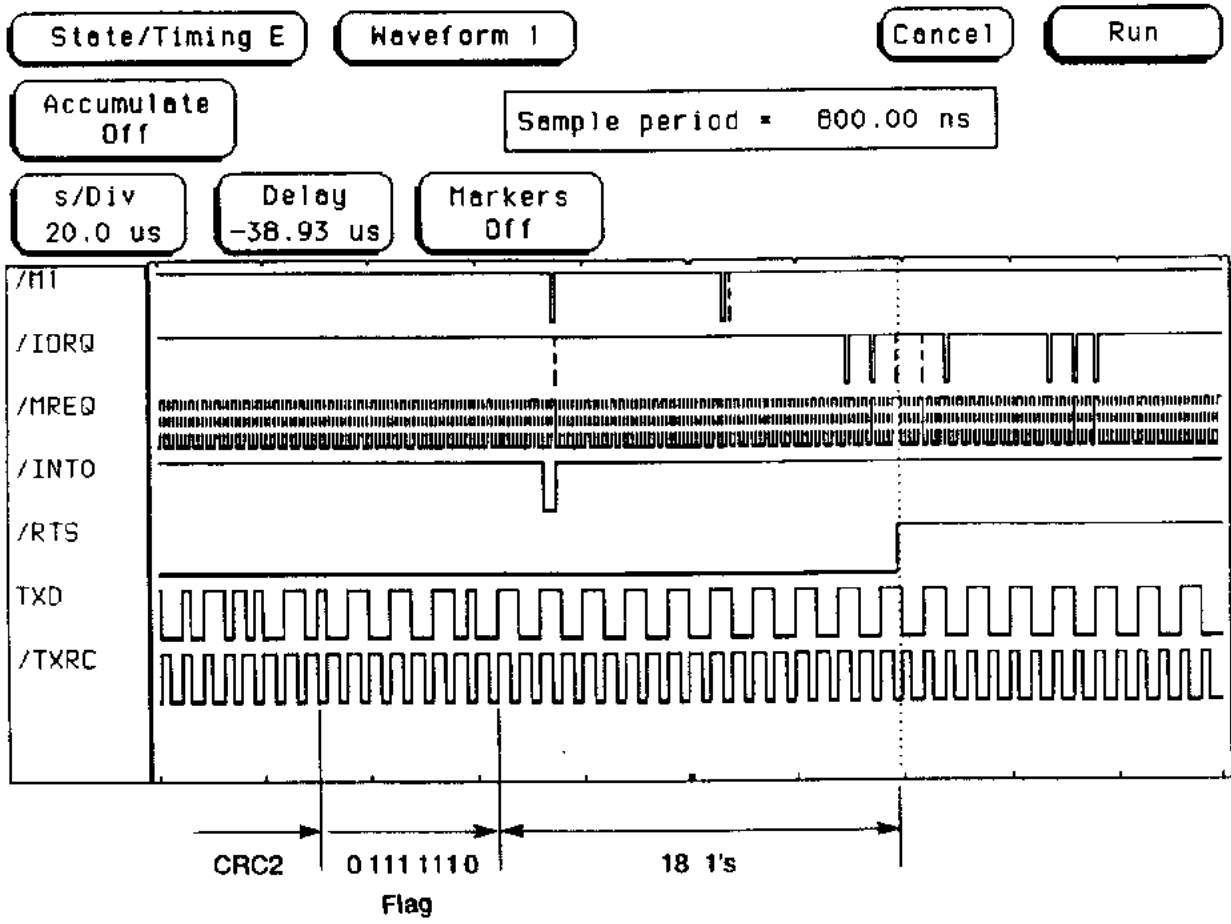


```

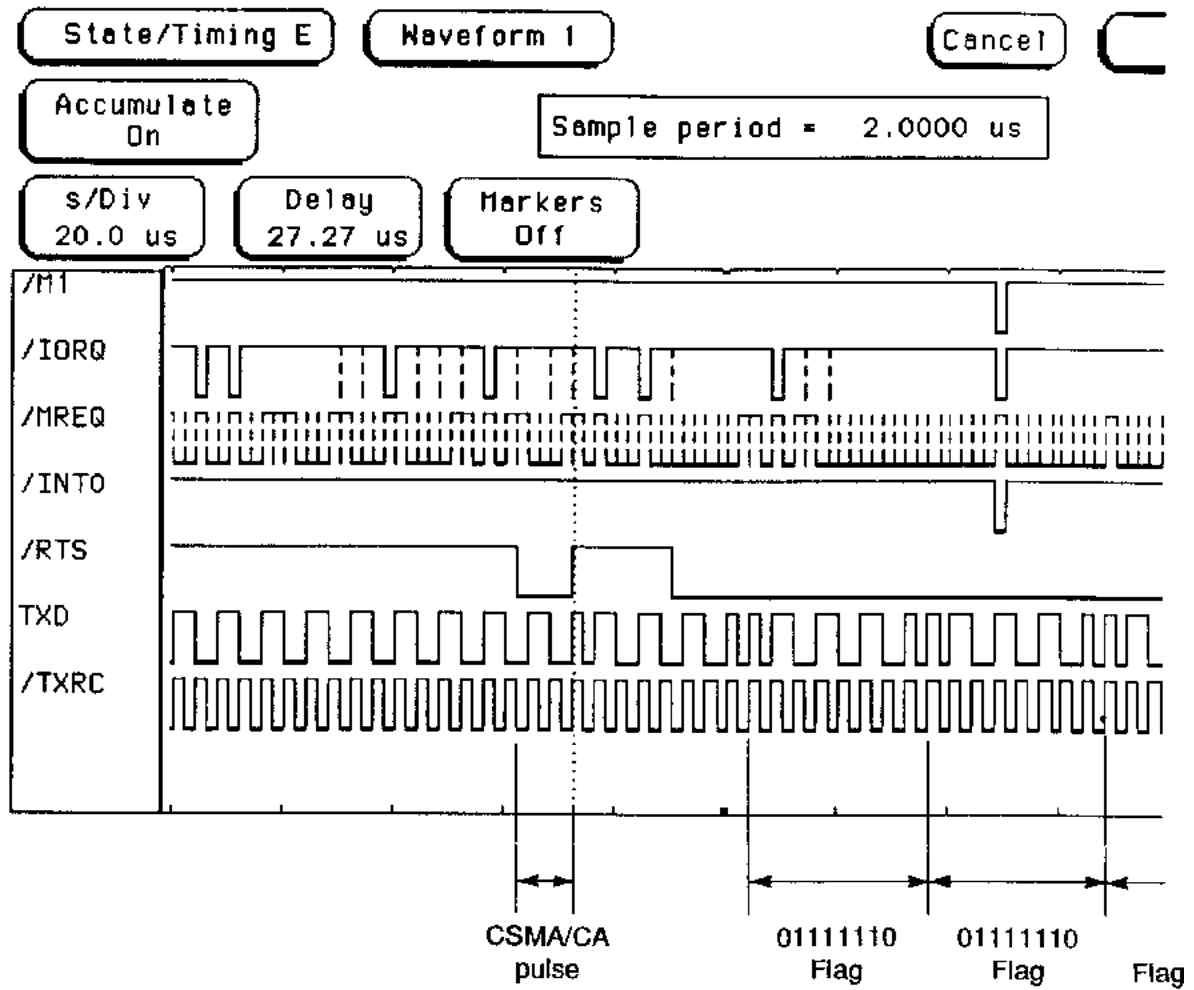
00001000          1365          org          1000h
00001000          1366 rx_buff: .block      length
1367
1388
1389              ;*****
1390              ;transmitter buffer area
1391              ;*****
0000b000          1392          org          0b000h
1398              ;
1399              ;*****
1400              ;transmit llap enq packet (3bytes)
1401              ;*****
0000b258 ff      1402 txlapenq:db          0ffh      ;broadcast id
0000b259          1403          .block      1          ;guess at myaddress
myaddress
0000b25a 81      1404          db          81h          ;llap enq type
1405              ;

```

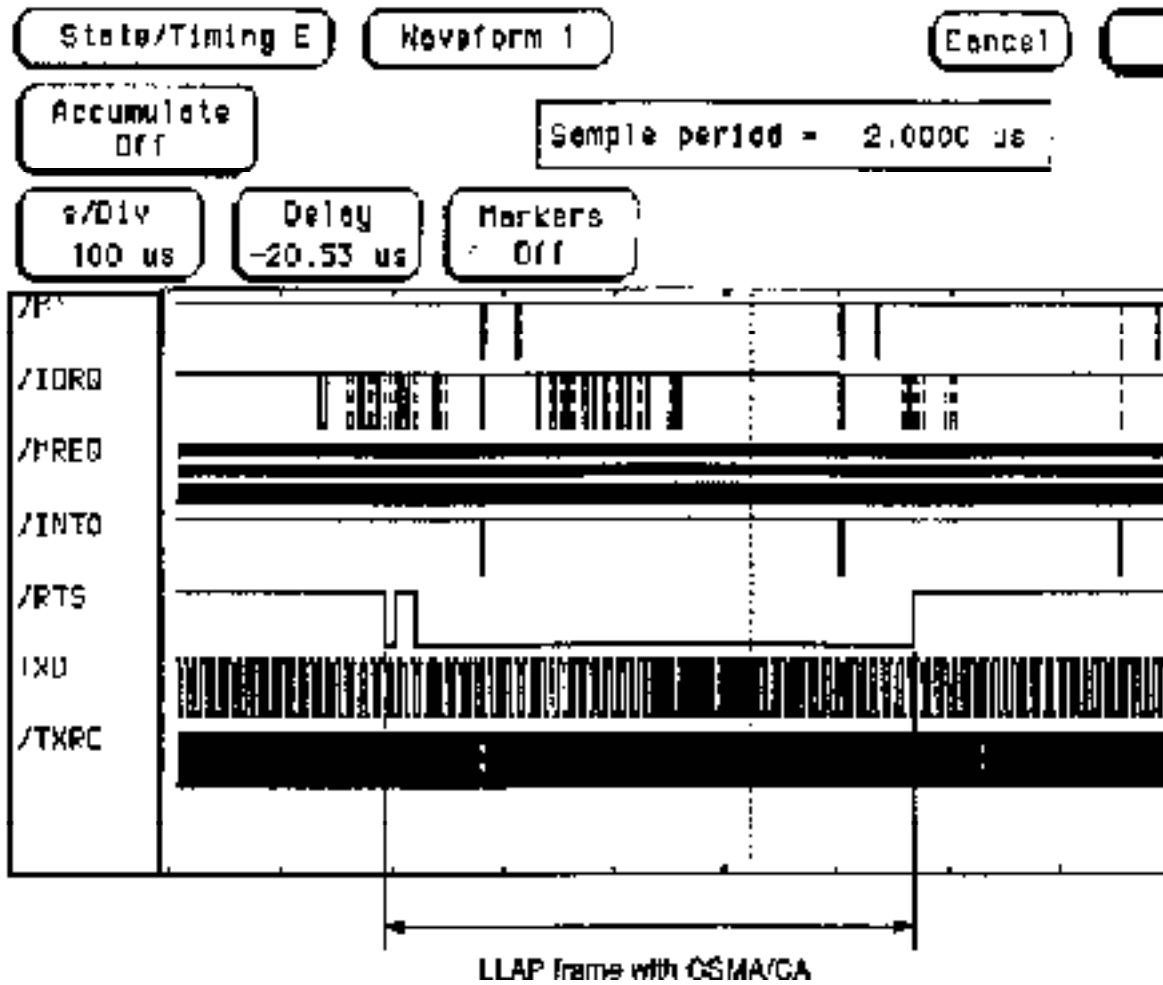
APPENDIX B



12 to 18 1's at the end of an LLAP frame



CSMA/CA before an LLAP frame



An LLAP Frame